# Characterisation of Parallel Independence in AGREE-Rewriting 

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# Contents 

Partial arrow classifier
AGREE-rewriting
Gluing construction
Residual
Parallel independence
Characterisation

## Partial Arrow Classifier

A category has partial arrow classifiers, if the following object-indexed family of morphisms exists:

For every object 0 , there is a monomorphism $\eta_{0}: 0 \rightarrow 0^{\circ}$ which satisfies the following universal property:

For every pair of morphisms ( $i: D \rightarrow X, f: D \rightarrow 0$ ) with monic $i$, there is a unique morphism $(i, f)^{\bullet}: X \rightarrow 0^{\bullet}$ such that the pair $(i, f)$ is pullback of the pair $\left(\eta_{0},(i, f)^{\bullet}\right)$.

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A category has partial arrow classifiers, if for every object $O$ there is object $O^{\bullet}$ and partial morphism $\varepsilon o: O^{\bullet} \rightarrow 0$ such that for every object $X$ and partial morphism $(p: X \rightarrow 0)$ there is unique total morphism $p^{\bullet}: X \rightarrow O^{\bullet}$ with $\varepsilon_{0} \circ \iota\left(p^{\bullet}\right)=p$, where functor $\iota$ is given by: $\iota: 0 \longmapsto 0$ and $\iota m: m \longmapsto[i d, m]$.

The embedding from category $C$ (with total arrows) into the category of partial arrows over $C$ is a free construction!

## Partial Arrow Classifier

Pushouts are hereditary
Pushouts preserve monomorphism
Pushouts along monomorphisms are pullbacks
Category has epi-mono-factorisation
Pullbacks are preserved by embedding ......
The embedding from category $C$ (with total arrows) into the category of partial arrows over $C$ is a free construction!

## Partial Arrow Classifier: Set



## Partial Arrow Classifier: Set

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## Partial Arrow Classifier: Set



# Partial Arrow Classifier: Graph 



## Partial Arrow Classifier: Graph

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# Partial Arrow Classifier: OO-Model 



## Partial Arrow Classifier: OO-Model

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## Partial Arrow Classifier: OO-Model


,Inversion‘ of Matches

,Inversion‘ of Matches

,Inversion‘ of Matches

,Inversion‘ of Matches


Pullback


## AGREE-Rewriting

## AGREE-Rewriting



## AGREE-Rewriting

Rule:


## AGREE-Rewriting

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## AGREE-Rewriting



## AGREE-Rewriting



## AGREE-Rewriting

Context Rule:
Inverse Male:
Base Match:
Trace:


## AGREE: Practical Example

Extract abstract type (Version I):


## AGREE: Practical Example

Extract abstract type (Version 2):


## AGREE: Local Copies



## AGREE: Local Copies



## AGREE: Local Copies



## AGREE: Global Copies



## AGREE: Global Copies



## AGREE: Local Deletion



## AGREE: Local Deletion



## AGREE: Global Deletion



## AGREE: Global Deletion



## AGREE: Local Addition



## AGREE: Local Addition



## Gluing Construction



## Gluing Construction

$$
\begin{aligned}
& (r, l) \circ(q, p)=(n, m) \circ(h, g) \\
& L \stackrel{I}{\longleftrightarrow} \mathrm{~K} \xrightarrow{r} R
\end{aligned}
$$

$$
\begin{aligned}
& \downarrow^{n} \text { (FPC) } \downarrow^{v} \text { (PO) } \downarrow q \\
& \mathrm{G} \underset{\mathrm{~g}}{\longleftarrow} \mathrm{D} \xrightarrow[h]{\longrightarrow} \mathrm{H}
\end{aligned}
$$

# Gluing for DPO-Rewriting 



## Gluing for DPO-Rewriting

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## Gluing for SPO-Rewriting



## Gluing for SqPO-Rewriting



# Gluing for AGREE-Rewriting 



# Gluing for AGREE-Rewriting 



## Gluing Construction



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## Gluing Construction



Gluing diagrams compose and decompose like pushouts

## Parallel Independence



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## Residual



## Residual



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## Residual



## Residual



# Characterising Independence 



## Characterising Independence

Match $m_{1}$ for rule I has residual after applying rule 2 at $\mathrm{m}_{2}$, only if
I. everything that $m$ I needs (locally copies, deletes, or preserves) is neither copied nor deleted (neither locally nor globally) by rule 2 at match $\mathrm{m}_{2}$.

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I. everything that $m$ I needs (locally copies, deletes, or preserves) is neither copied nor deleted (neither locally nor globally) by rule 2 at match $\mathrm{m}_{2}$.
2. everything that rule I adds is neither (globally) copied nor deleted by rule 2 at match $\mathrm{m}_{2}$.

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## Characterising Independence



## Conclusion

AGREE-rewriting is instance of the Gluing Construction!
There is a precise notion of residual!
Gluing and mutual residuals provides Church-Rosser!
Residuals can be characterized syntactically!

Are global effects useful?

## Thank you for your attention

