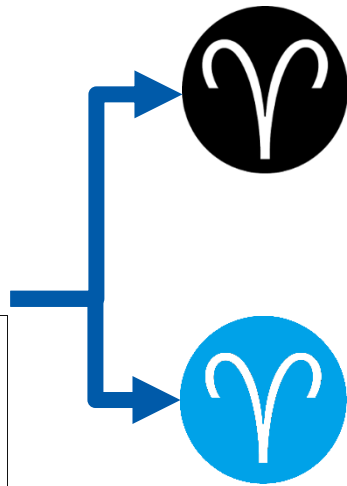
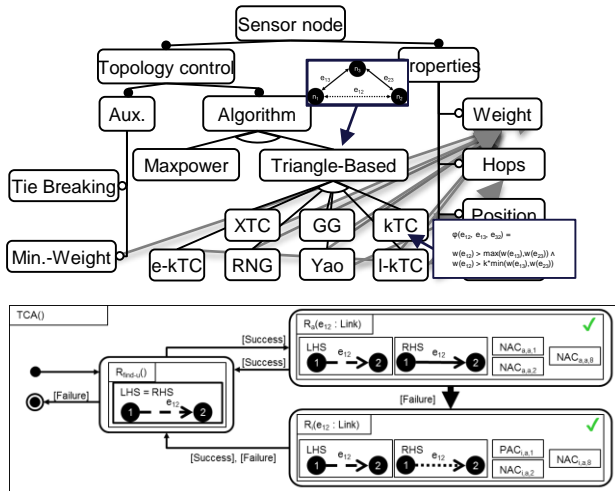


A systematic approach to constructing families of incremental topology control algorithms using graph transformation

Presentation at ICGT 2018
Toulouse, France, 2018-06-25



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Supported by the Cooperative Research Center 1053 "MAKI" – <https://tiny.cc/MAKI>

Technische Universität Darmstadt
Fachgebiet Echtzeitsysteme – Real-Time Systems Lab

Prof. Dr. rer. nat. Andy Schürr

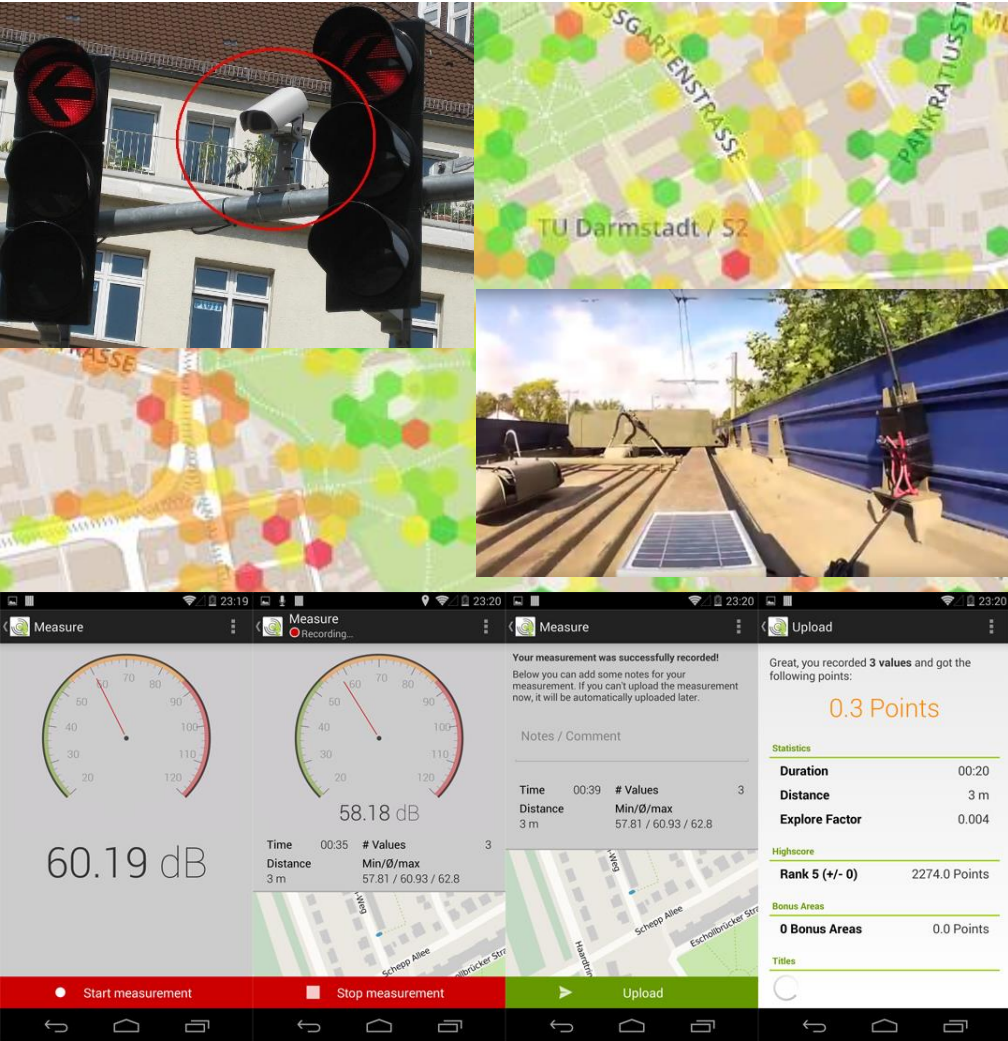
Dept. of Electrical Engineering and Information Technology

Dept. of Computer Science (adjunct Professor)

[SoSyM17] Kluge, R. et al. SoSyM 2017, DOI: 10.1007/s10270-017-0587-8

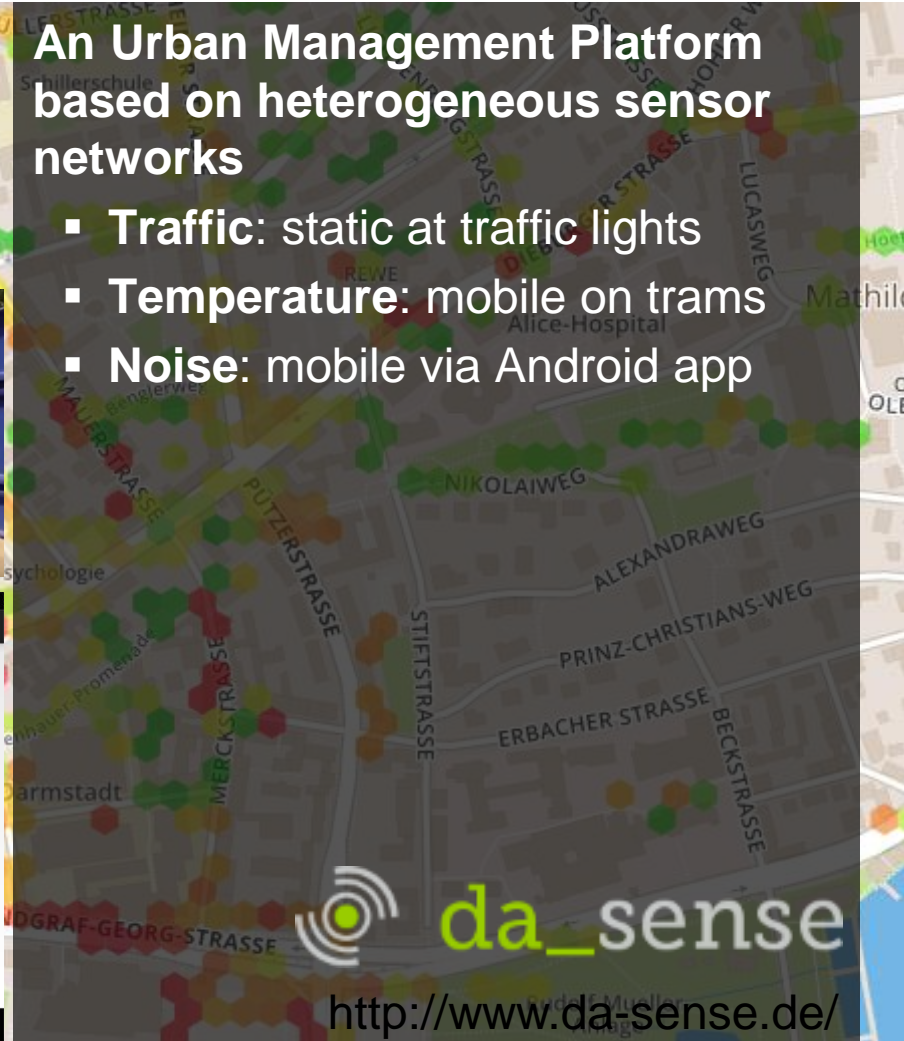
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Example: da_sense – A hybrid sensor network for Smart Cities (“Digitalstadt Darmstadt”) (I)

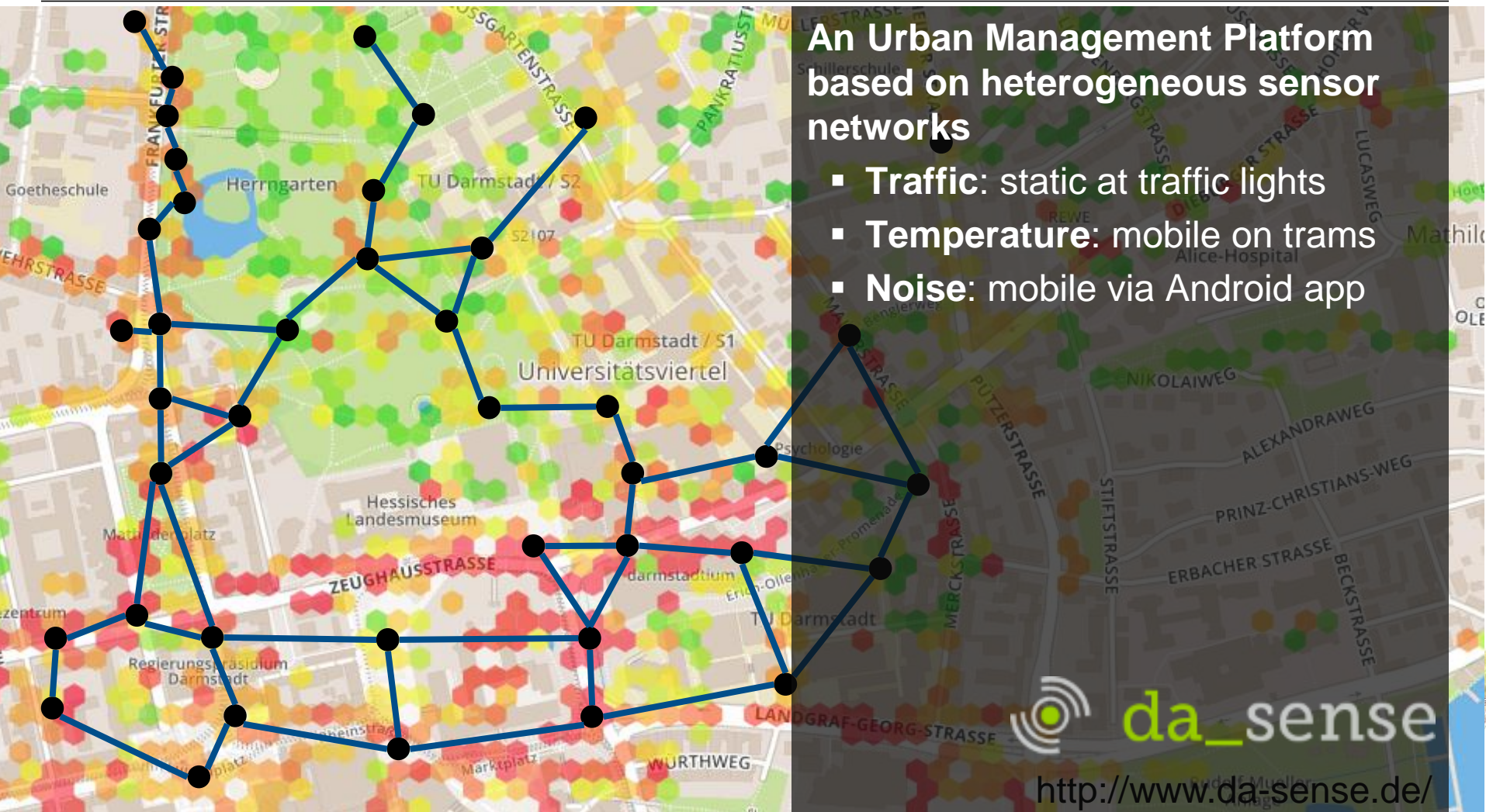


An Urban Management Platform based on heterogeneous sensor networks

- Traffic: static at traffic lights
- Temperature: mobile on trams
- Noise: mobile via Android app

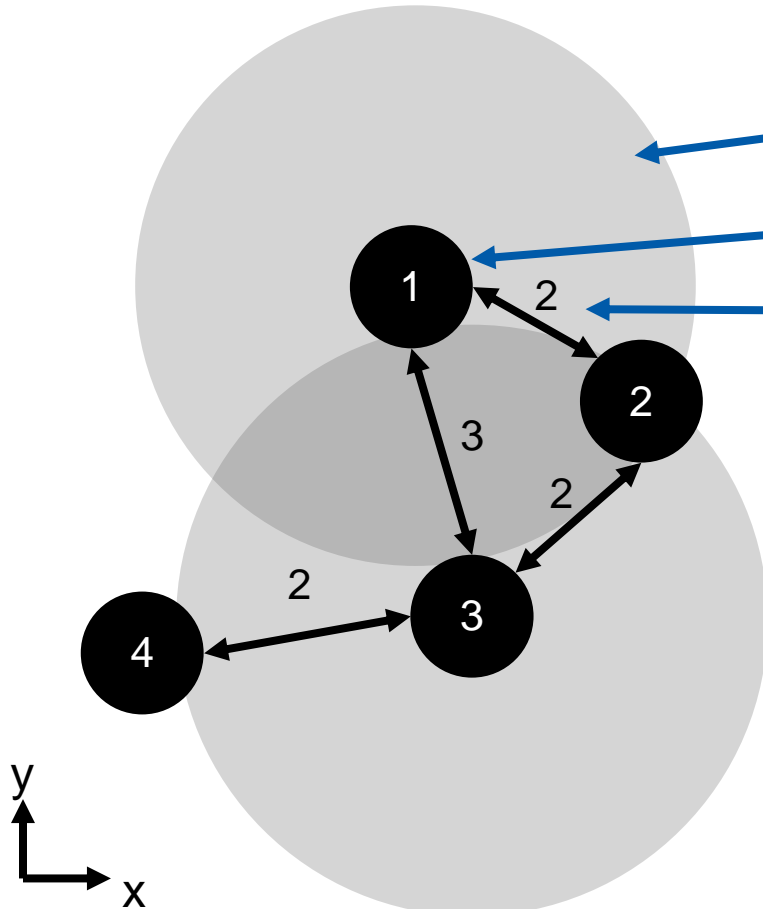


Example: da_sense – A hybrid sensor network for Smart Cities (“Digitalstadt Darmstadt”) (II)



Graph-based topology model for Wireless Sensor Networks

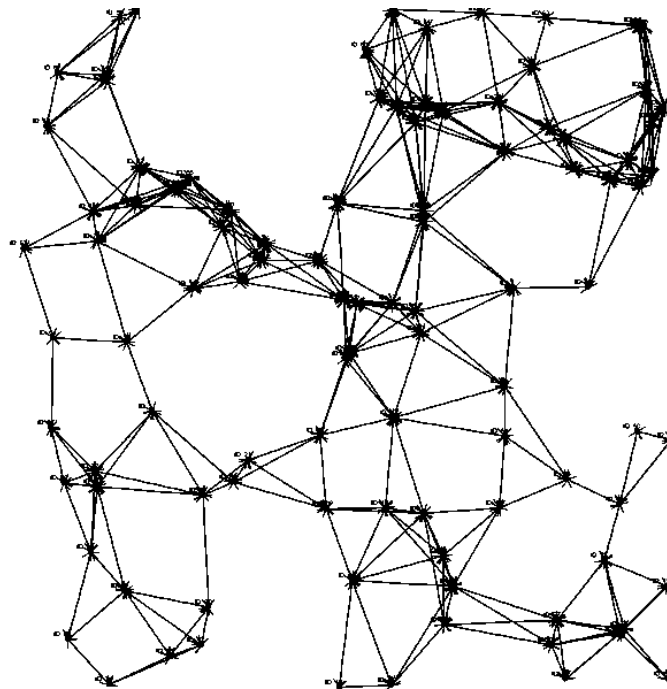
Wireless Sensor Network topology



TelosB
sensor node
48kB ROM,
10kB RAM,
2xAAA battery

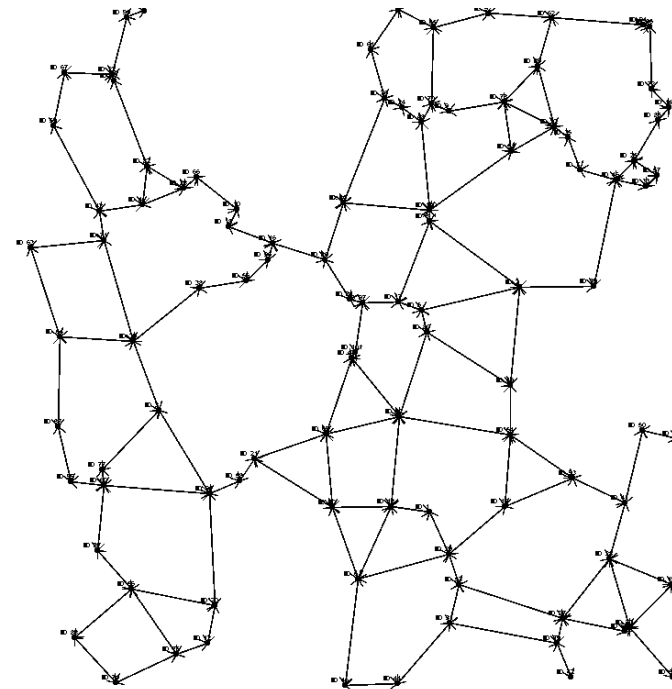


Topology control sparsens topologies to improve non-functional property



Input topology

Topo. Ctrl.



Output topology



Maintain
connectivity

— Communication link

● Sensor node



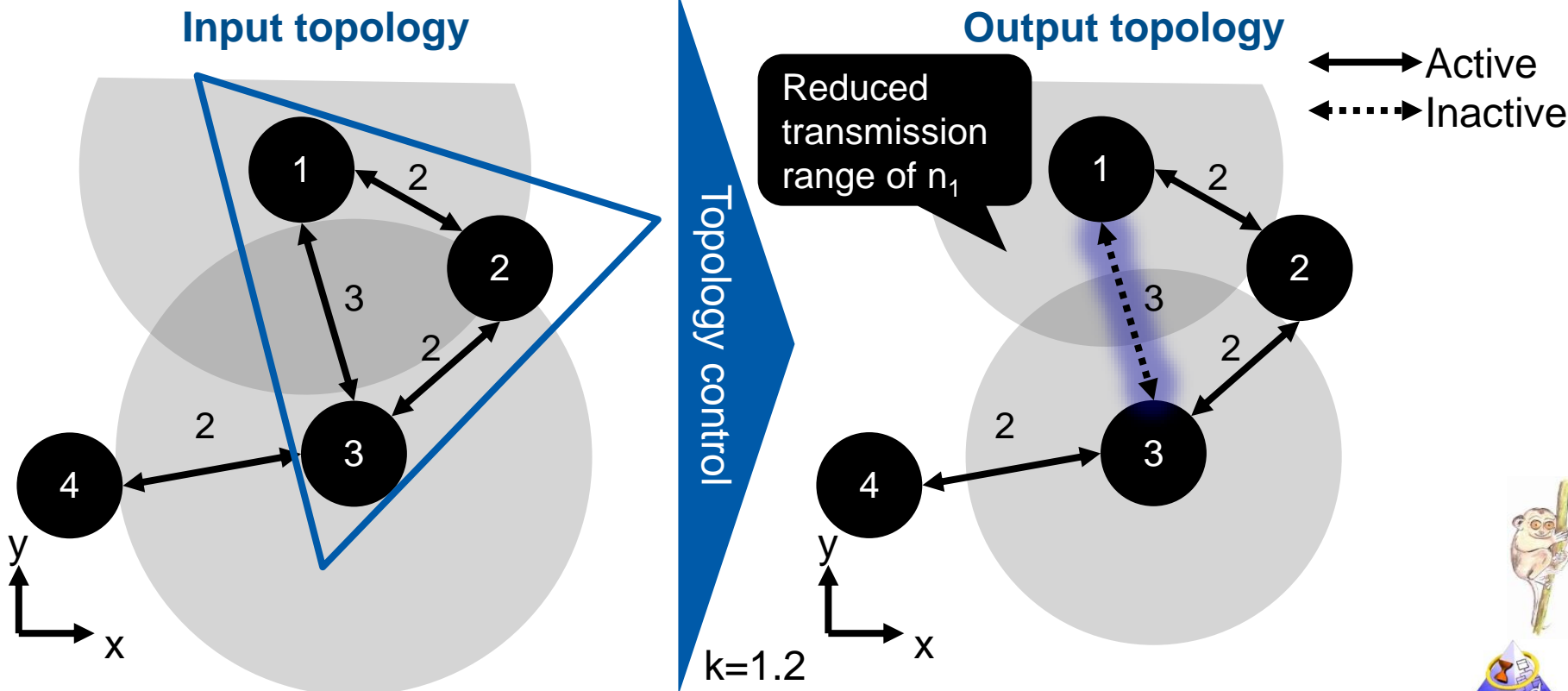
Topology control by example: kTC algorithm

kTC rule:

"Inactivate a link if (and only if) it is

(i) the weight-maximal link in a triangle and

(ii) at least k -times longer than the weight-minimal link in the triangle."



The curse of low abstraction in traditional communication system development

Theorem V.1. $G_{KTC} \subseteq G_{GG}$, or equivalently, the diametric circle of any two nodes $u, v \in G_{KTC}$ is empty.

Proof. We will show that $(u, v) \notin G_{GG}$ implies that $(u, v) \notin G_{KTC}$. Pick a $(u, v) \in G - G_{GG}$. Then there must exist a $w \in G$ such that w lies inside the diametric circle of u and v . By the assumption of the UDG, $(u, v) \in G$ implies that $(u, w) \in G$ and $(v, w) \in G$. Without loss of generality, assume u and v are oriented horizontally. The maximum value of $\min(d(u, w), d(v, w))$ is then attained on the top or bottom of the the diametric circle where $d(u, w) = d(v, w)$.

Theorem V.1. $G_{KTC} \subseteq G_{GG}$, or equivalently, the diametric circle of any two nodes $u, v \in G_{KTC}$ is empty.

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Corollary V.2. G_{KTC} is planar.

Proof: The Gabriel Graph G_{GG} is planar [7]. □

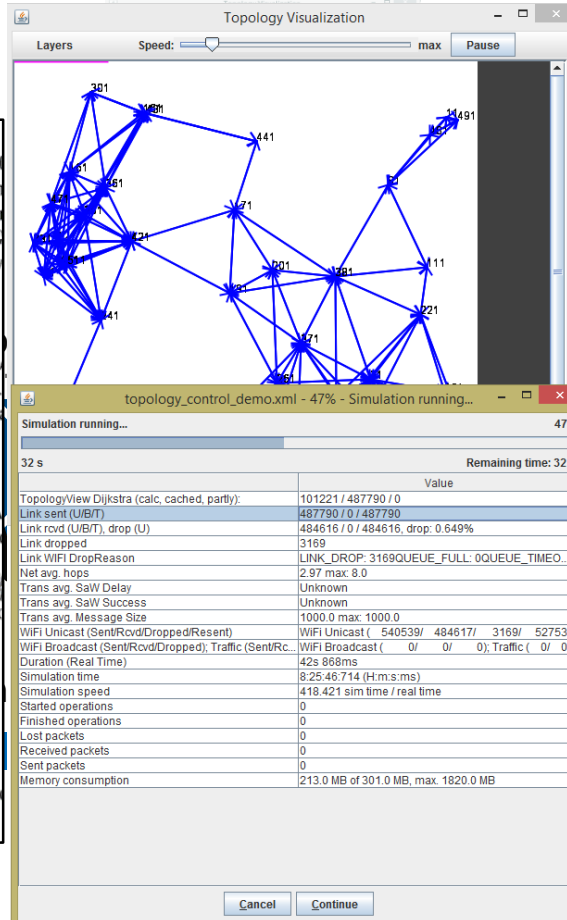
Theorem V.3. $G_{KTC} = G_{XTC}$ whenever $k = 1$.

Proof: For clarity, we ignore the tie-breaking case, where both XTC and KTC discard the same edge based on IDs.

In XTC an edge (u, v) is removed iff there is a node w with $d(u, w) < d(u, v)$ and $d(v, w) < d(u, v)$. Nodes u, v , and w form a triangle where (u, v) is the longest edge. When $k = 1$ KTC removes exactly the longest edge. □

Corollary V.4. $G_{XTC} \subseteq G_{KTC}$

Proof: Increasing k only adds edges. □

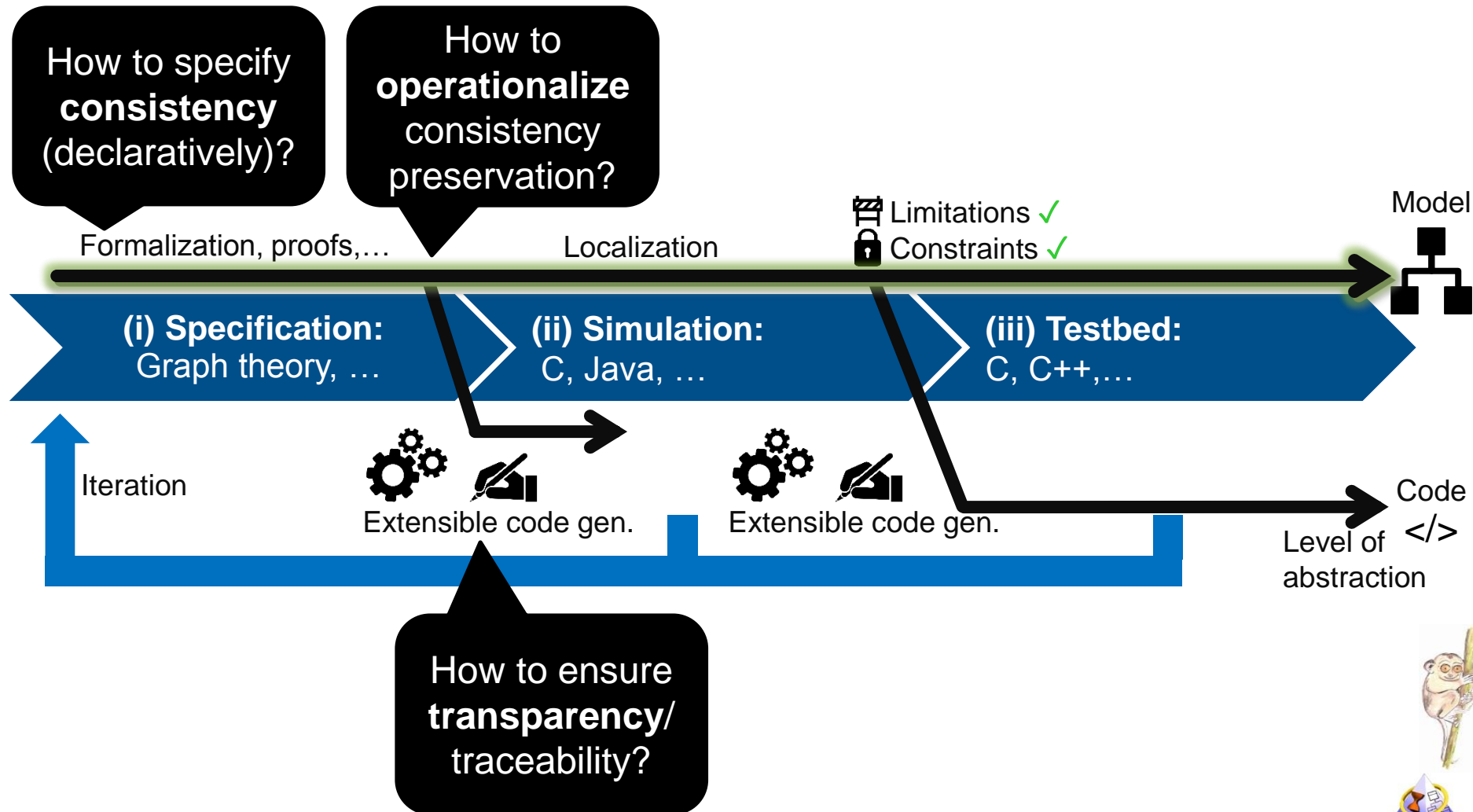


The screenshot shows a 'Topology Visualization' window with a network graph and a 'topology_control_demo.xml - 47% - Simulation running...' window with a table of metrics.

Simulation running... 47%	
Remaining time: 32 s	
	Value
Topology/View/Dijkstra (calc, cached, partly):	101221 / 487790 / 0
Link sent (UB/T):	487790 / 0 / 487790
Link road (U/B/T), drop (U):	484616 / 0 / 484616, drop: 0.649%
Link dropped:	3169
Link WiFi DropReason:	LINK_DROP: 3169 / QUEUE_FULL: 0 / QUEUE_TIMEOUT: 0
Net avg. hops:	2.97 max: 8.0
Trans avg. SaW Delay:	Unknown
Trans avg. SaW Success:	Unknown
Trans avg. Message Size:	1000.0 max: 1000.0
WiFi Unicast (Sent/Rcvd/Dropped/Resent):	WiFi Unicast (540539 / 484617 / 3169 / 52753)
WiFi Broadcast (Sent/Rcvd/Dropped): Traffic (Sent/Rcvd/Dropped):	WiFi Broadcast (0 / 0 / 0); Traffic (0 / 0 / 0)
Duration (Real Time)	42s 868ms
Simulation time	8:25:46.714 (H:m:s.ms)
Simulation speed	418.421 sim time / real time
Started operations	0
Finished operations	0
Lost packets	0
Received packets	0
Sent packets	0
Memory consumption	213.0 MB of 301.0 MB, max: 1820.0 MB



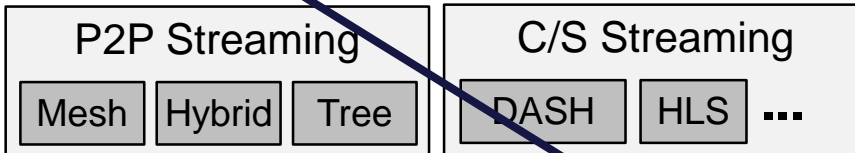
Goal: Support the correct-by-construction development of topology control algorithms



MAKI for a better Future Internet

Multi-Mechanismen-Adaption für das Künftige Internet

C: Communication systems (concrete self-adaptive systems)



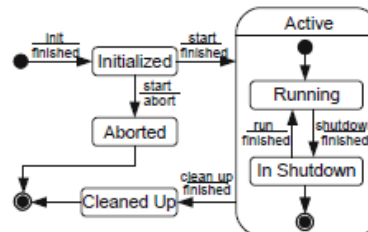
B: Adaptation mechanisms (reusable components for adaptive systems)



A: Construction methods (models, design patterns, languages)

Our area

- Network topologies and topology adaptation
- Specification languages
- Software engineering



```

1 filter(
2   max(
3     join(
4       match(TP, T1, self <- e0 -> n1,
5         n1 <- e1 -> n2),
6       match(TP, T3, self <- e2 -> n3),
7       e0.weight),
8     count(
9       match(TP, T2, self - e3 -> n4)) = 0)
10 execute every match:
11   at(self, TP, T2) add_neighbor(n1)

```

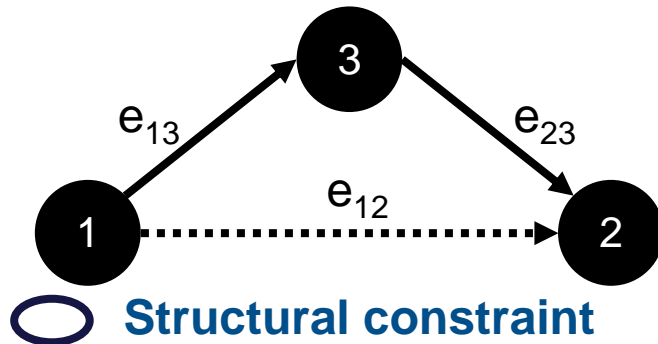


Unstructured consistency specification

Unstructured formulation

kTC: "Inactivate a link if (and only if) it is
 (i) the **weight-maximal link in a triangle** and
 (ii) at **least k-times longer than the weight-minimal link in the triangle**"

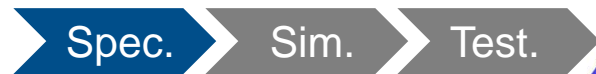
Structured formulation



$$\varphi(e_{12}, e_{13}, e_{23}) =$$

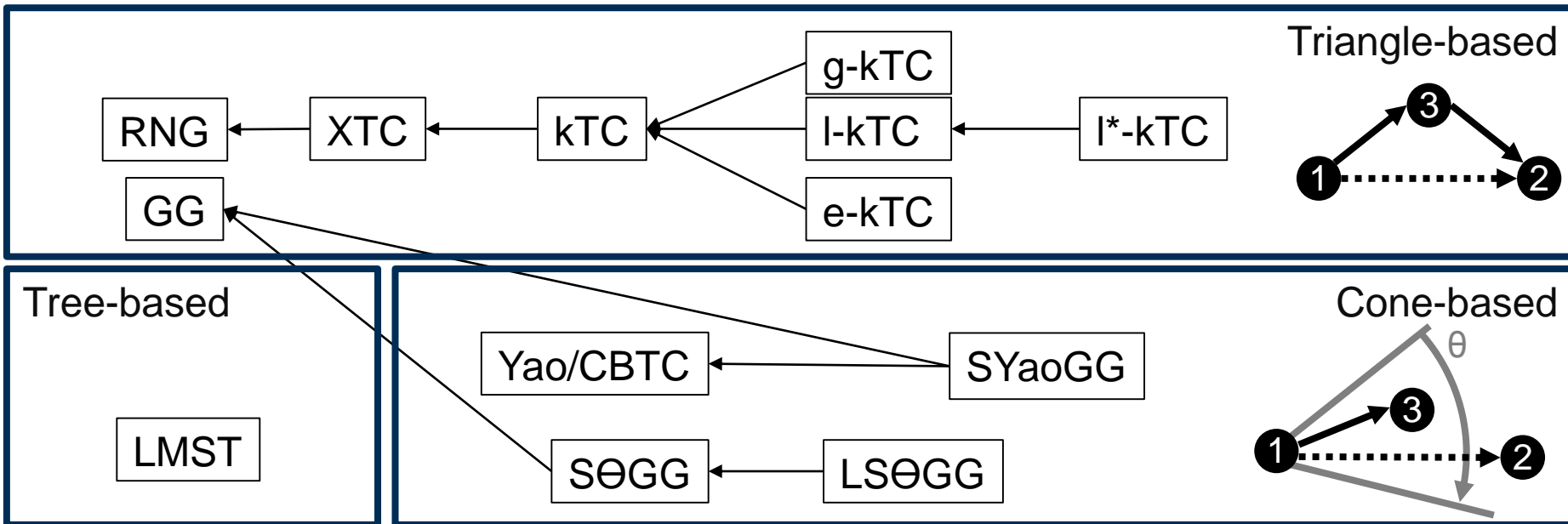
$$\wedge \quad w(e_{12}) > \max(w(e_{13}), w(e_{23})) \wedge \\ w(e_{12}) > k * \min(w(e_{13}), w(e_{23}))$$

Problem 1: Implicit, unstructured, or informal specification of constraints.



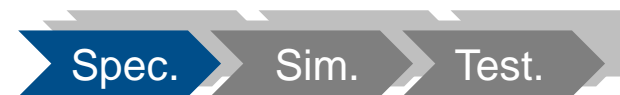
Topology control algorithms form families

- Family: common structural pattern
- Algorithm: refinement based on attribute constraints






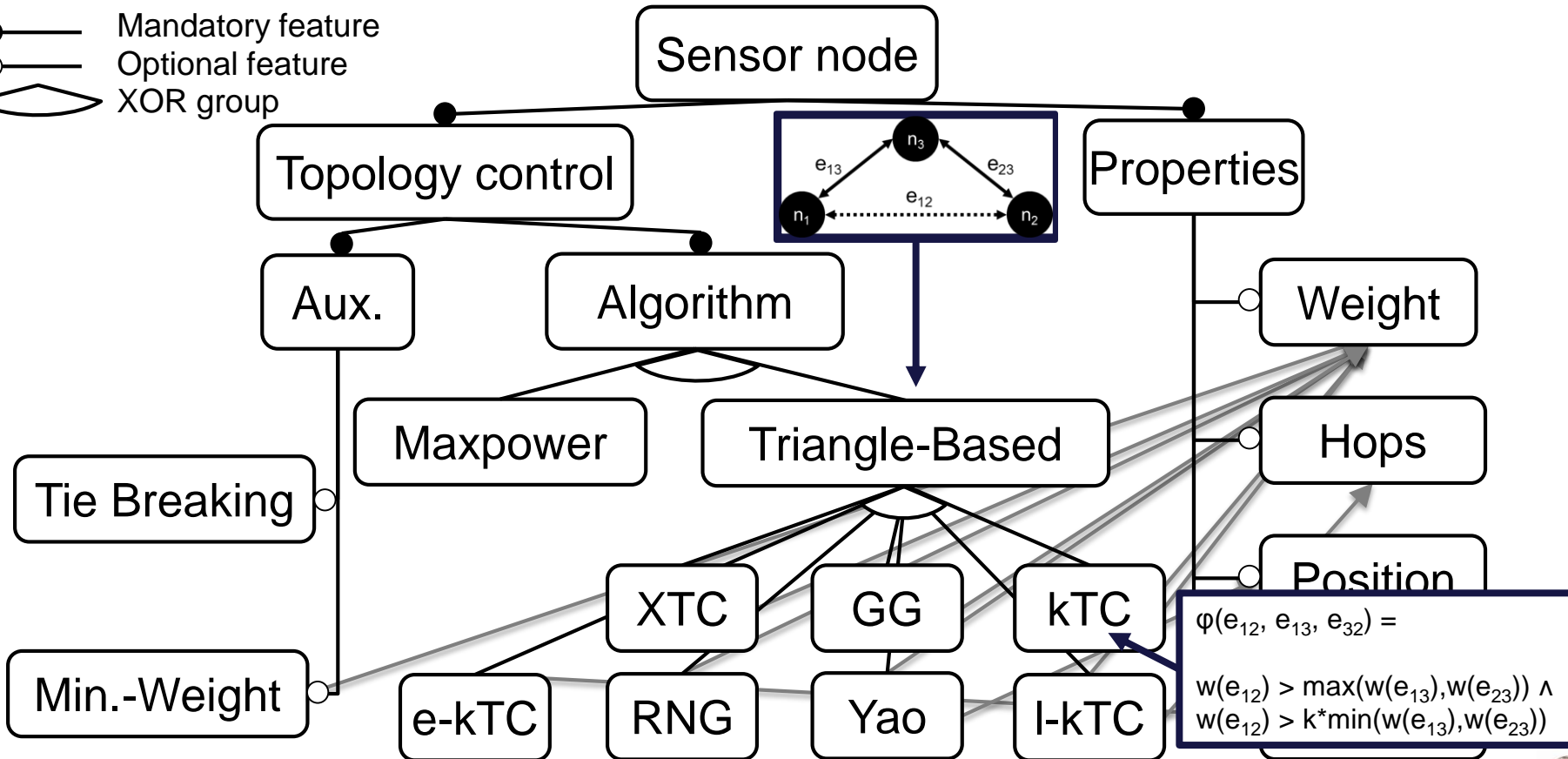
Problem 2: Insufficient usage of relationships among topology control algorithms

← builds on (uses/refines)



TC Algorithm families are Dynamic Software Product Lines

-  Mandatory feature
-  Optional feature
-  XOR group



Advantages: Reuse in of predicates, spec. of topology control reconfiguration

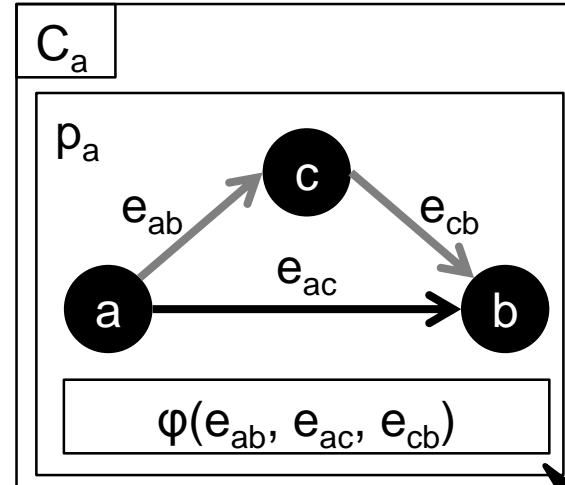
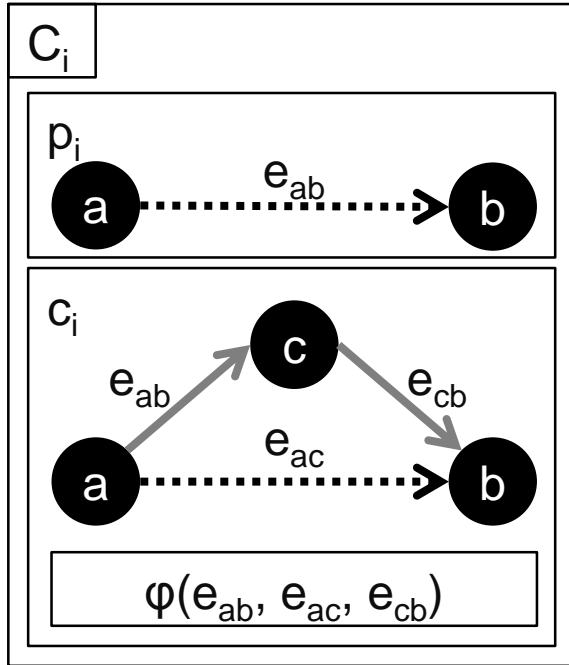
Spec.

Sim.

Test.



Graph constraints for specifying local consistency properties



 Active
 Inactive
 Don't care

Violated if match of p_a exists.

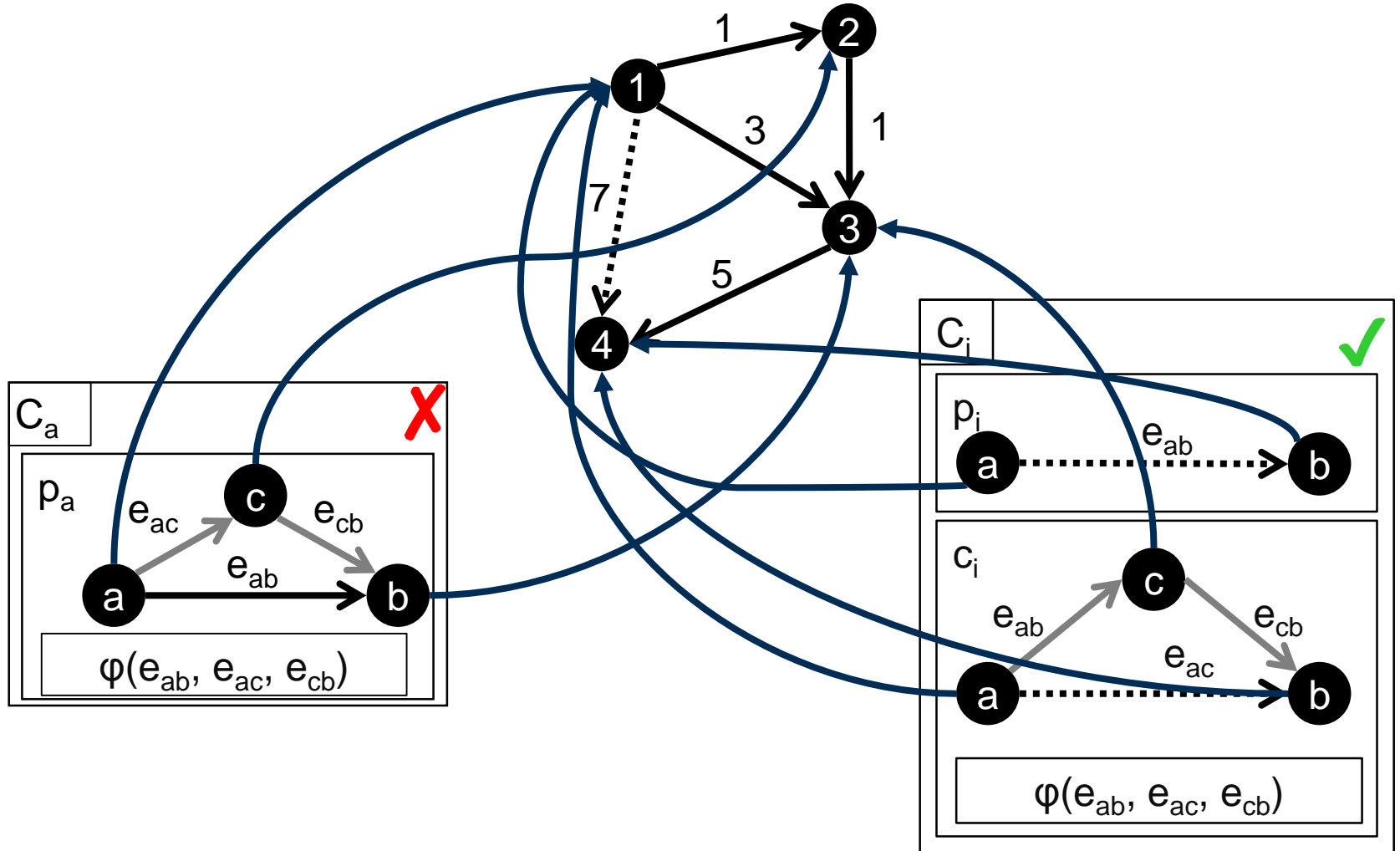
"Each inactive link should be part of a triangle for which φ holds"

"No active link should be part of a triangle for which φ holds"

Advantages: expressiveness, formal + domain-specific + operationalizable

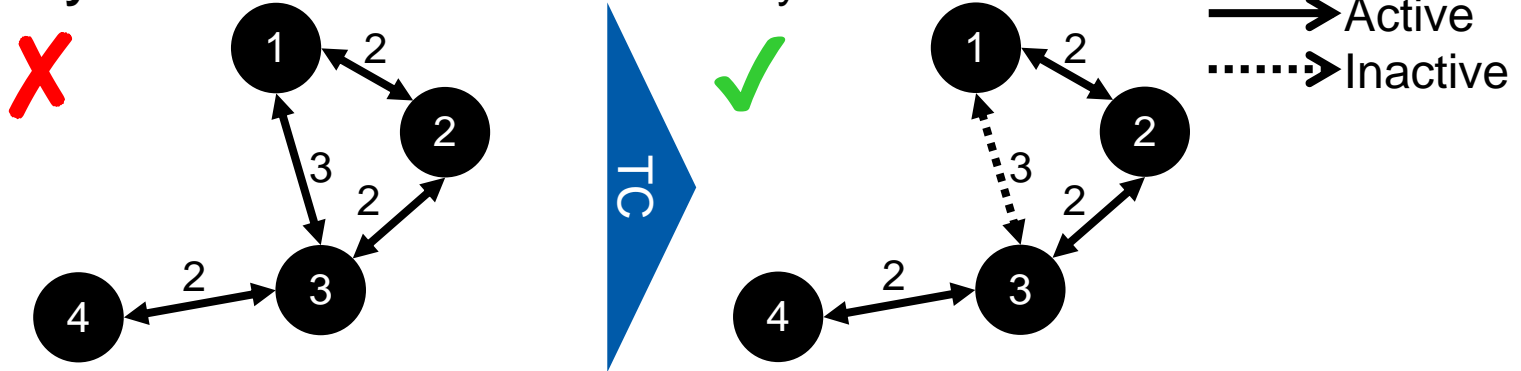


Example: Fulfilled and violated constraints

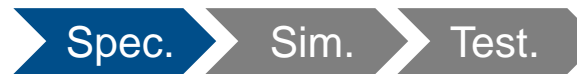
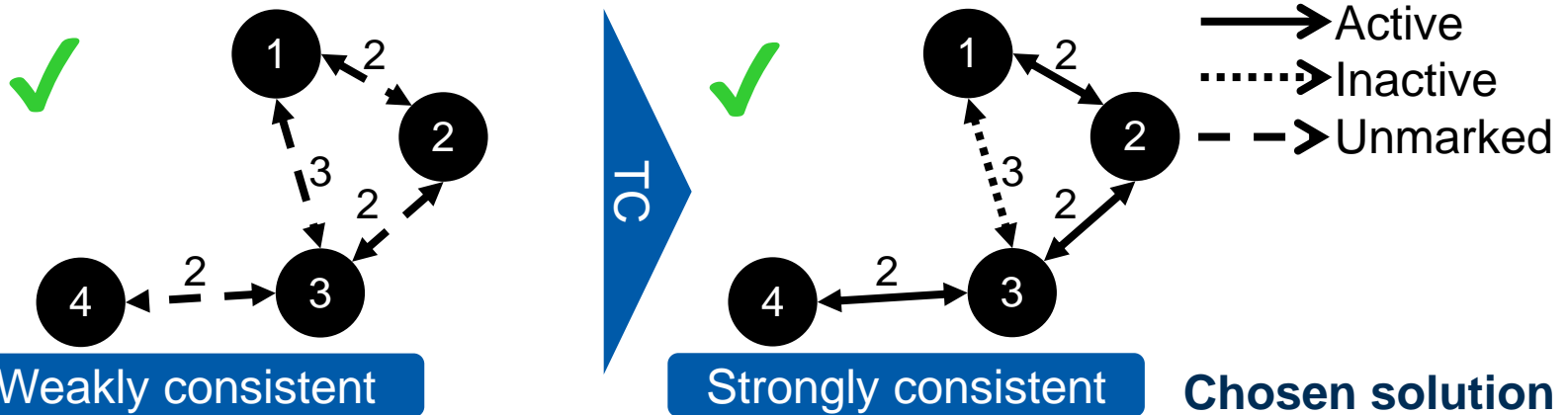


Enforcing and preserving consistency

Consistency enforcement: Recover from any inconsistent state



Consistency preservation: Weak and strong consistency



Recap on consistency preservation

Global consistency properties

implies

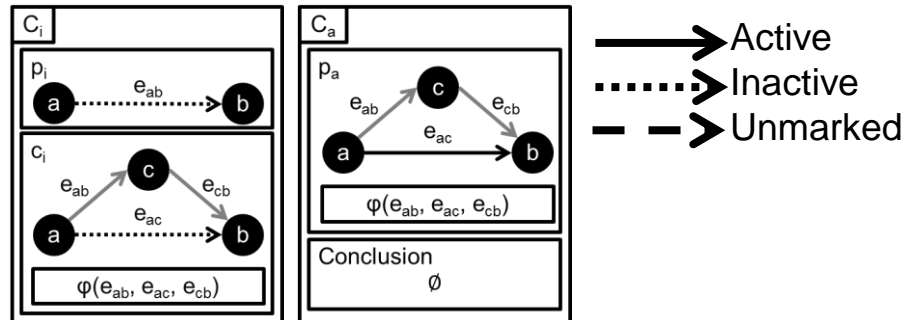
Local consistency properties (declarative)

preserves

Topology control algorithm specification (operationalized)

Generic
Algorithm-specific

→ e.g., connectivity
→ e.g., "no ϕ -triangles"



?

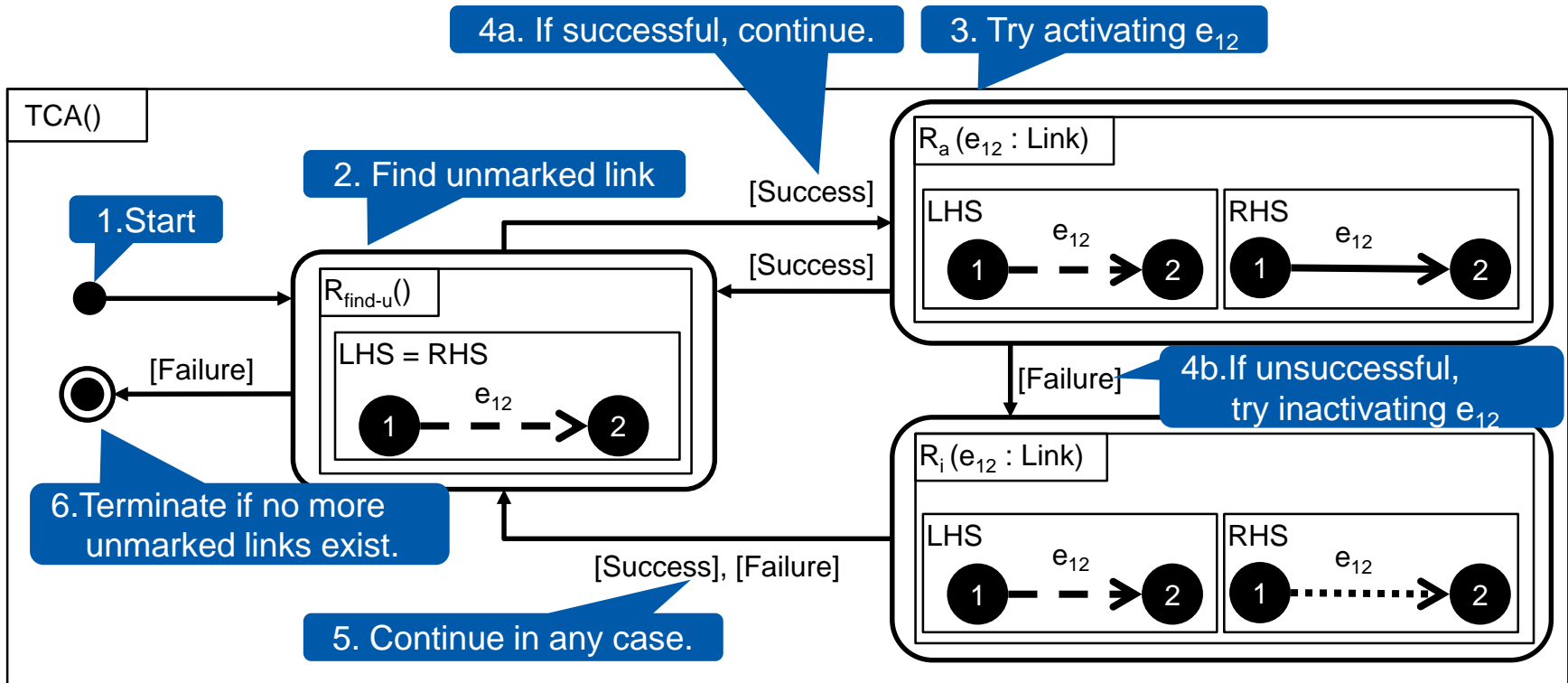
Spec.

Sim.

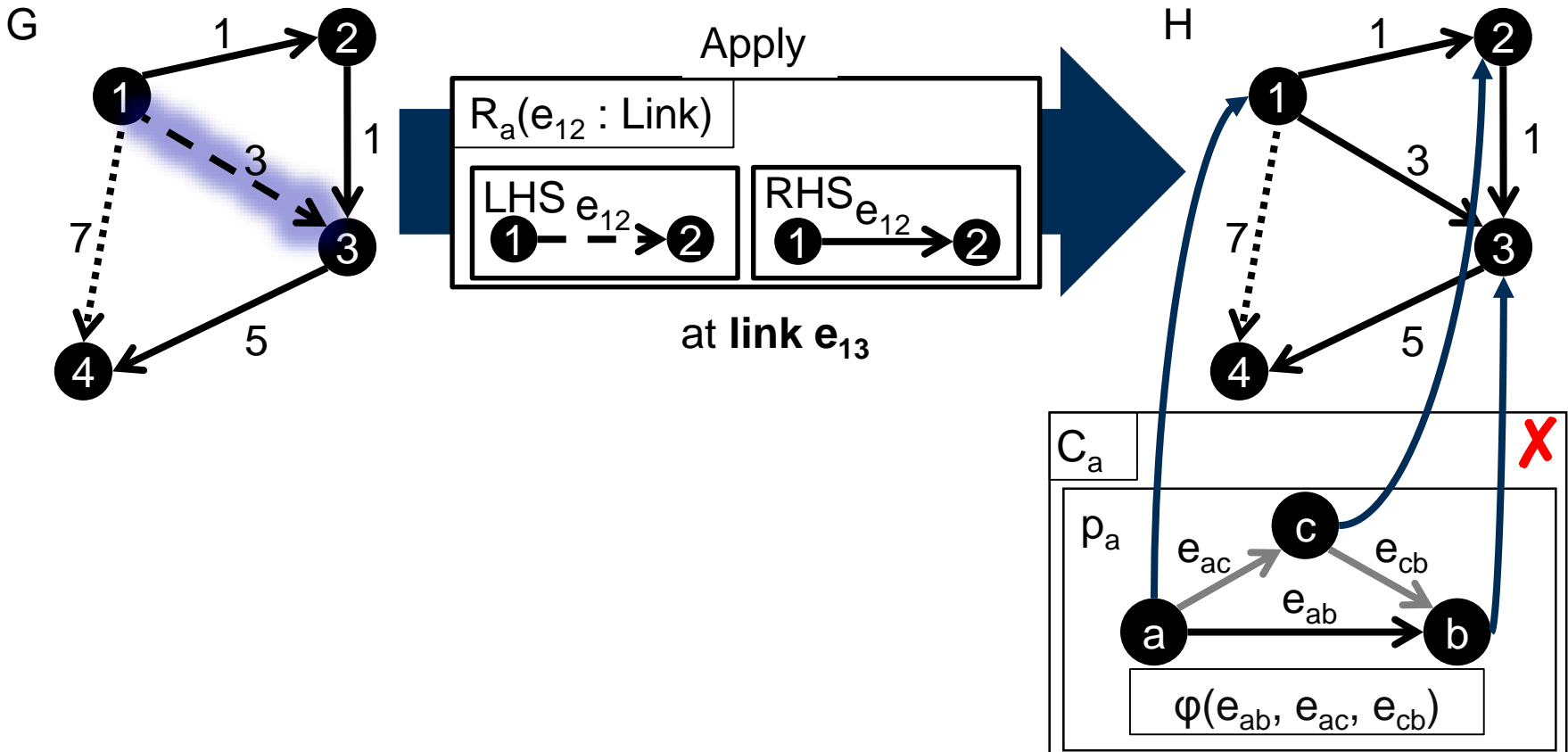
Test.



Specifying algorithmic implementation using Story-Driven Modeling [FNT+98]

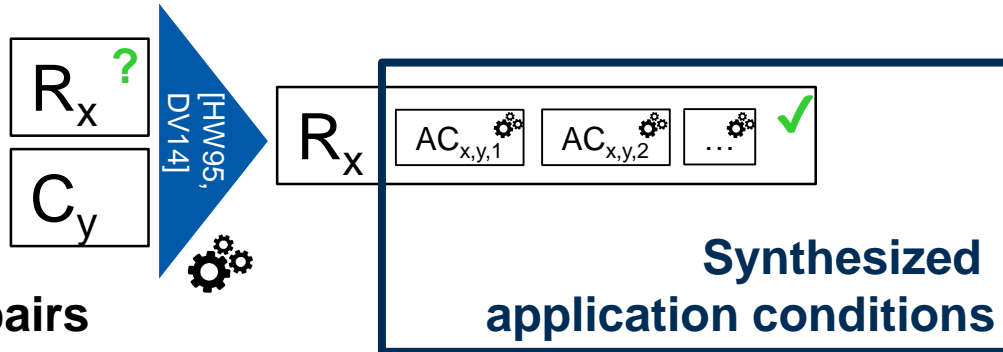


Rule applications may violate consistency!

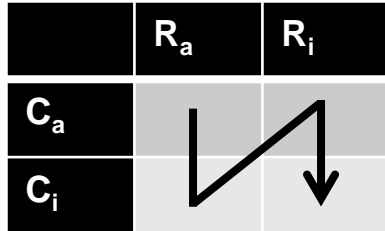


Ensuring inductive consistency preservation

Pairwise refinement

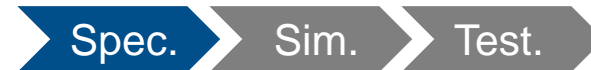
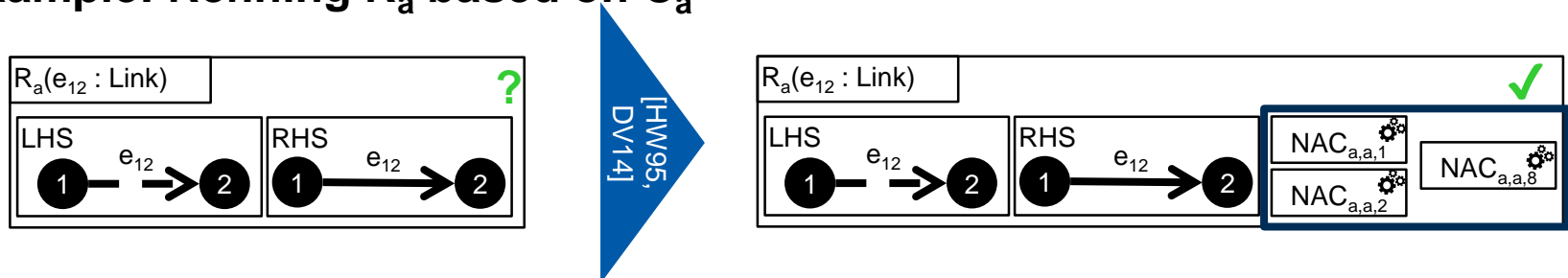


Repeat for all (R_x, C_y) pairs

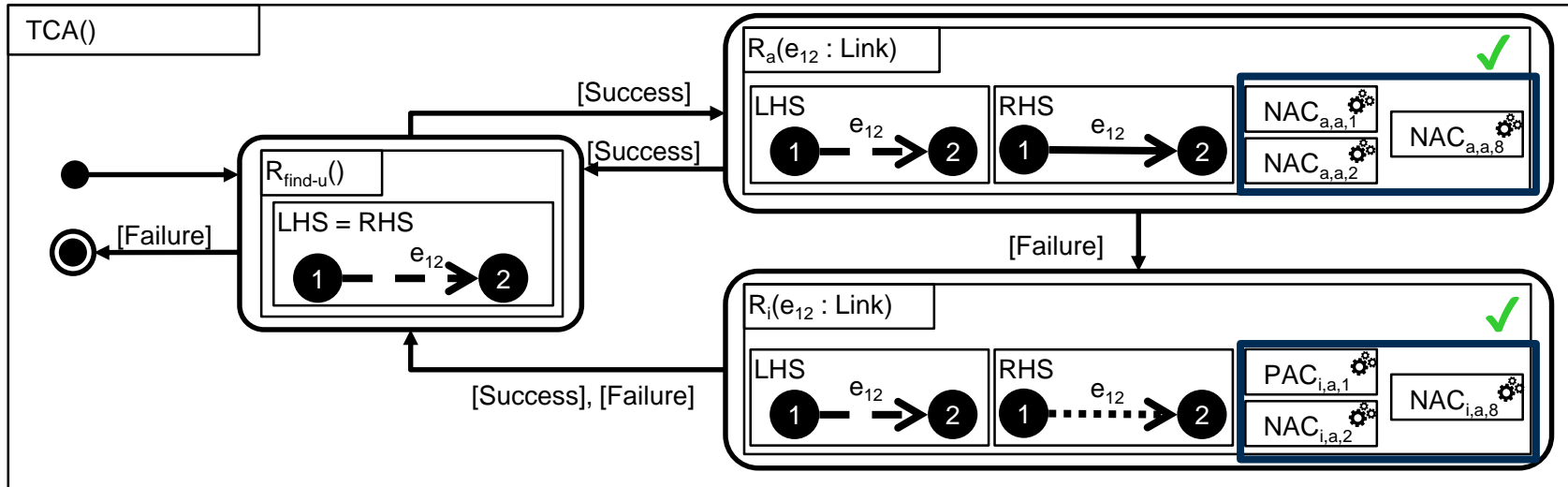


Advantages:
mechanical algorithm,
existing tool support

Example: Refining R_a based on C_a



Refined algorithm specification



Advantages: inductive invariant, strong consistency on termination

Required: proof of termination

Spec.

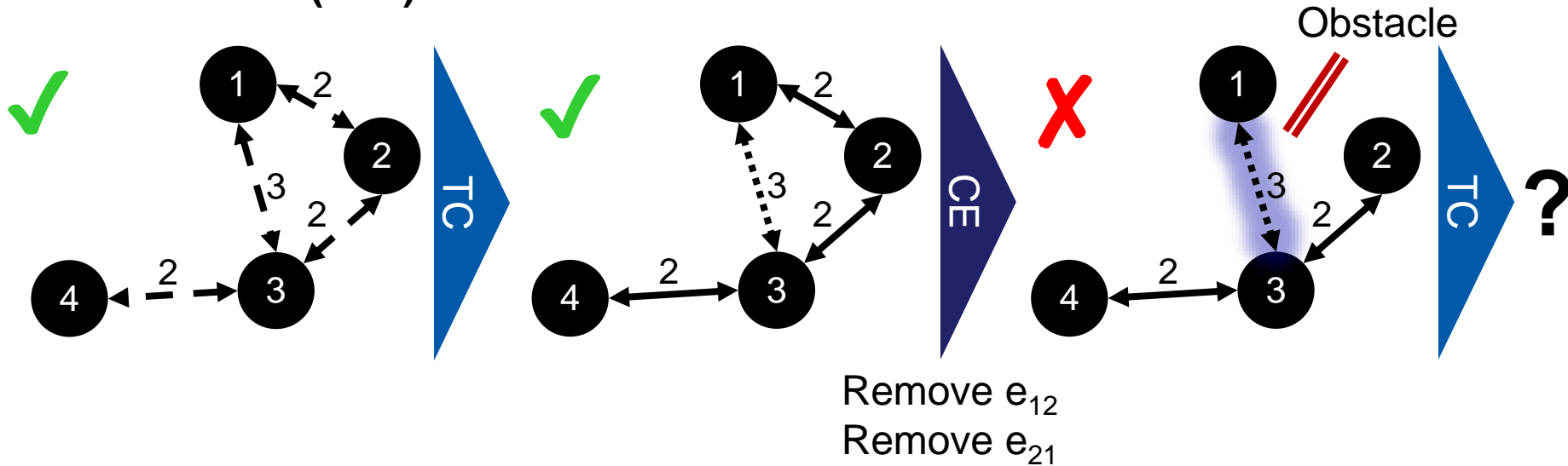
Sim.

Test.

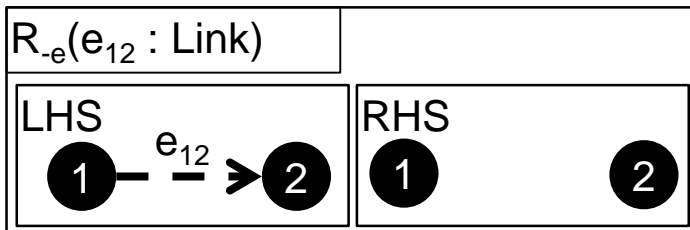


Topologies are never stable: Context events and dynamic topology control

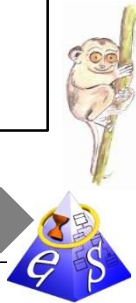
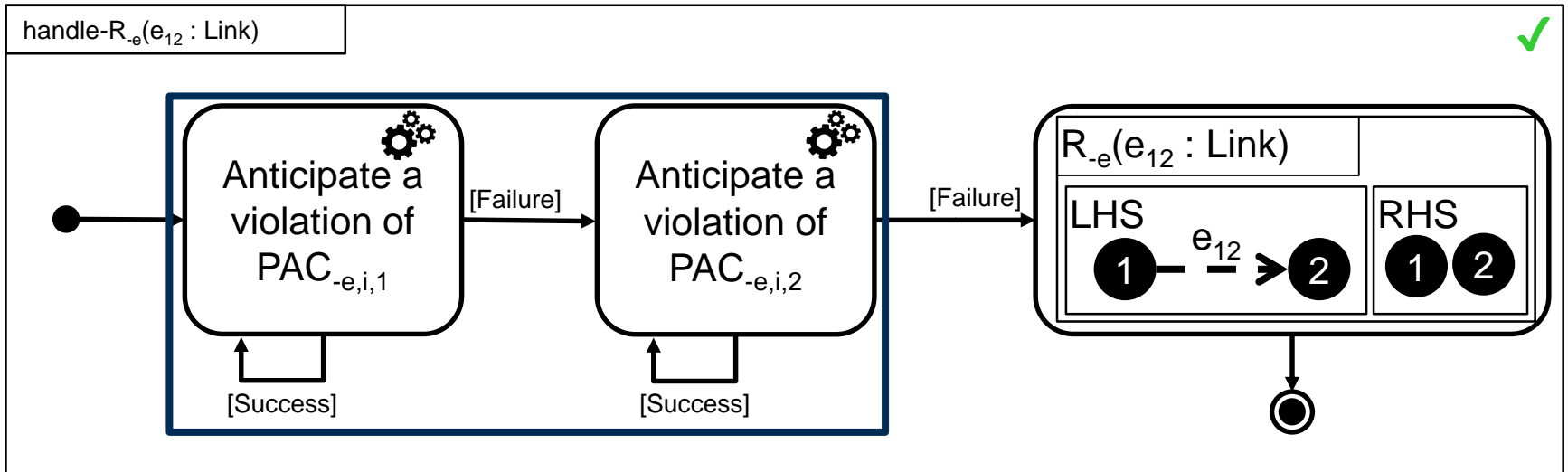
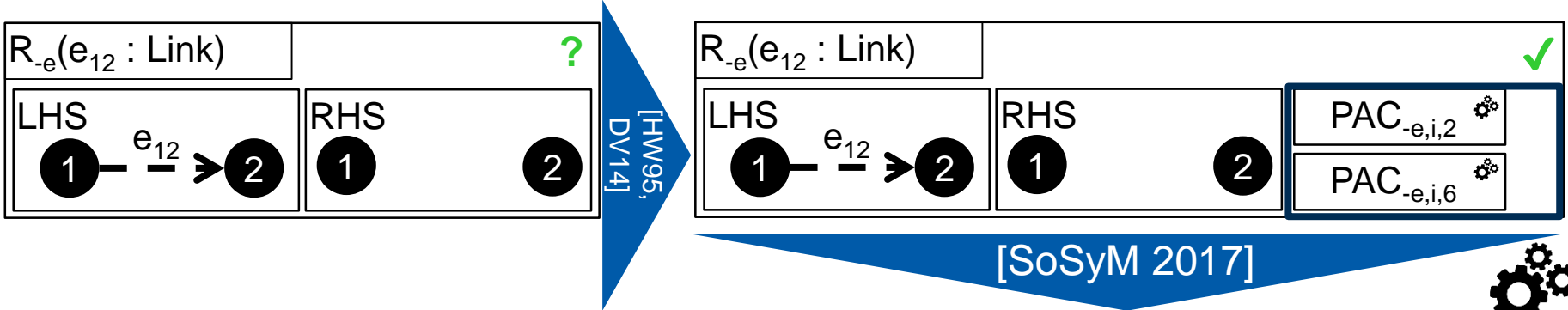
Context events (CEs) reflect environmental influences



Example: Link removal rule

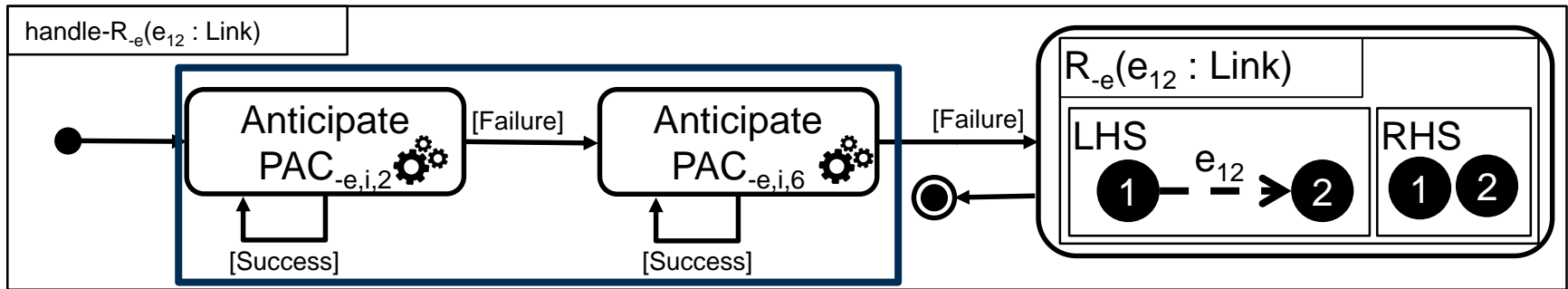


Context event handlers: Anticipating consistency violation

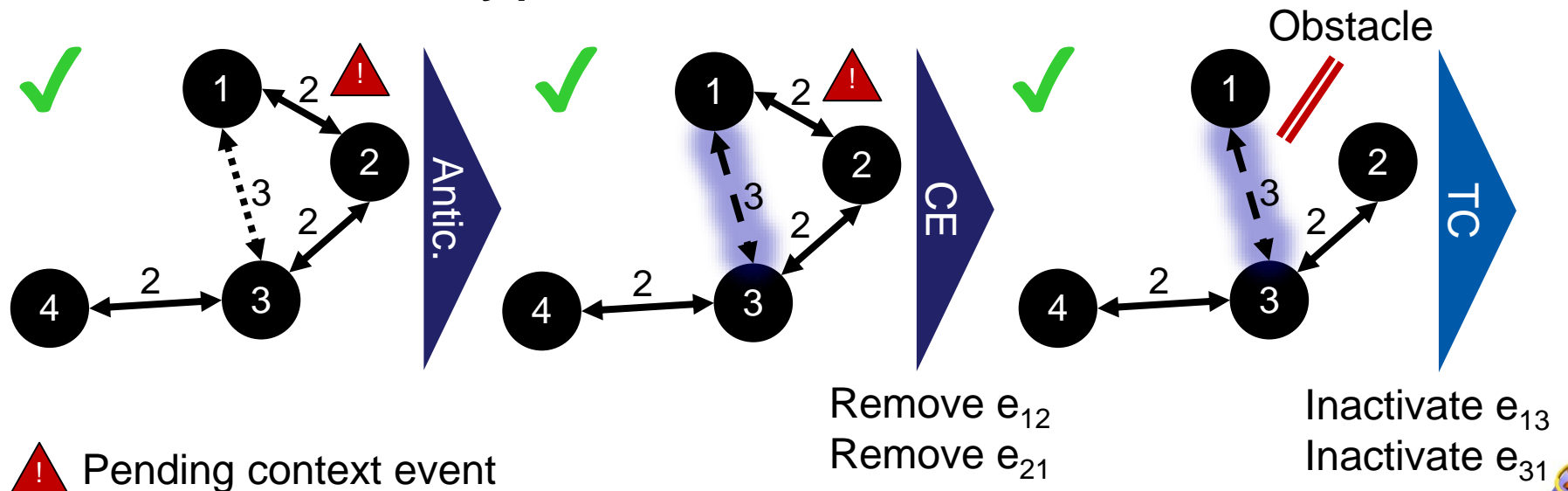


Consistency preservation by context event handling

– Context event handler for R_{-e}



– Achieved consistency preservation



Recap on specification phase

Global consistency
properties

implies

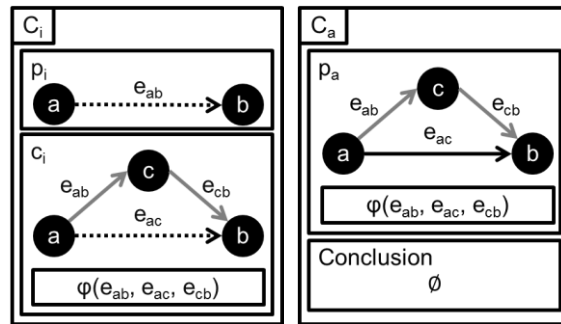
Local consistency
properties
(declarative)

preserves

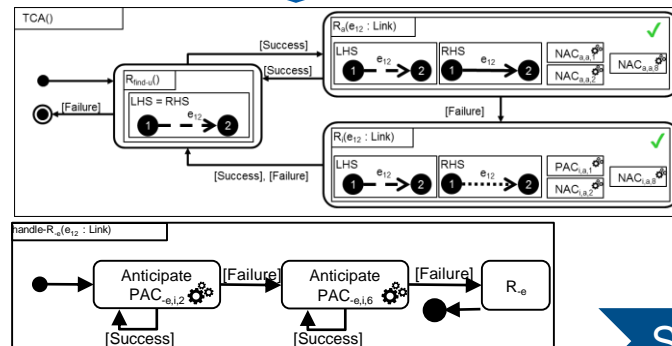
Topology control
algorithm specification
(operationalized)

Generic
Algorithm-specific

→ e.g., connectivity
→ e.g., "no ϕ -triangles"



[HW95, DV14]



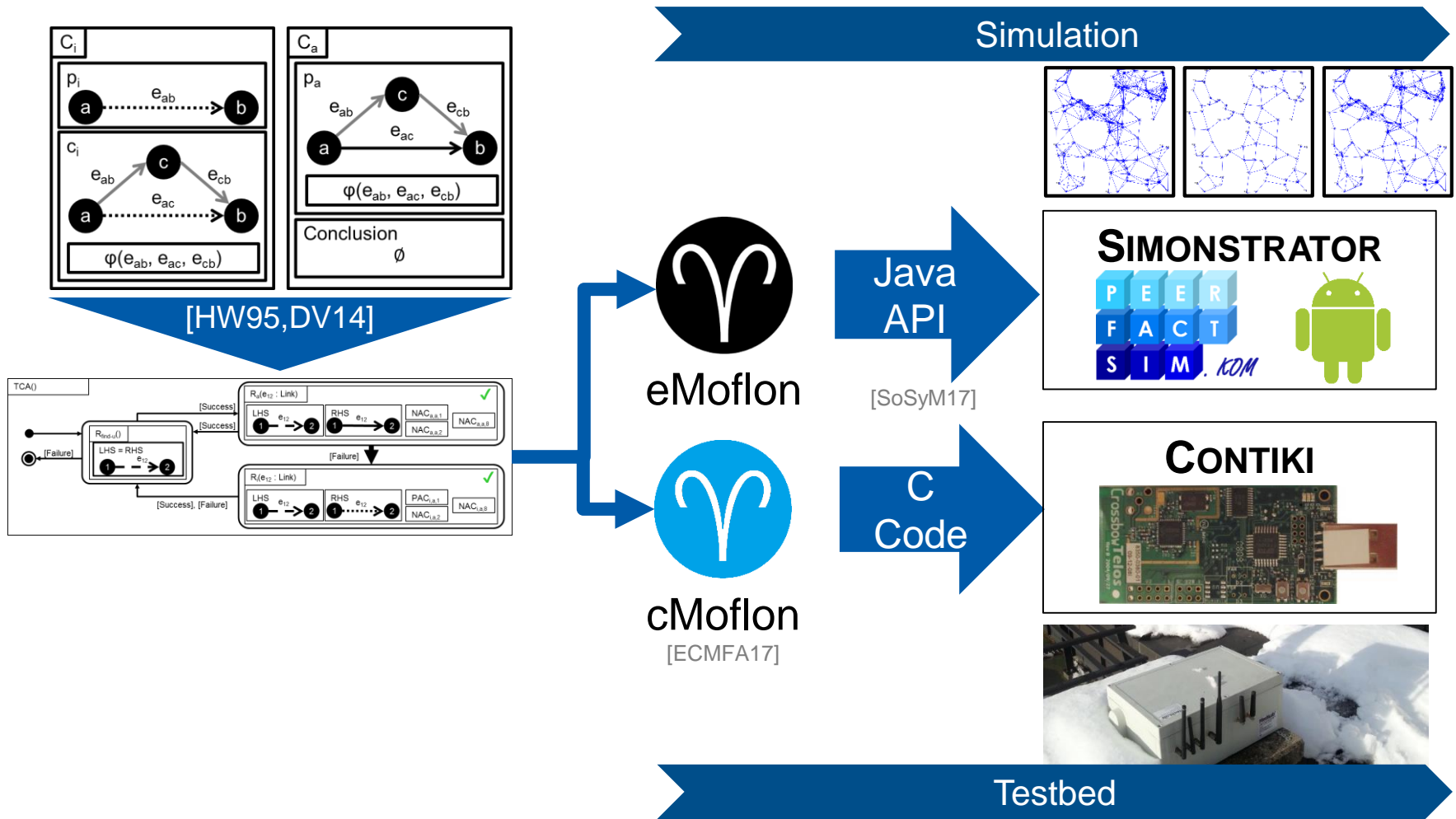
Spec.

Sim.

Test.



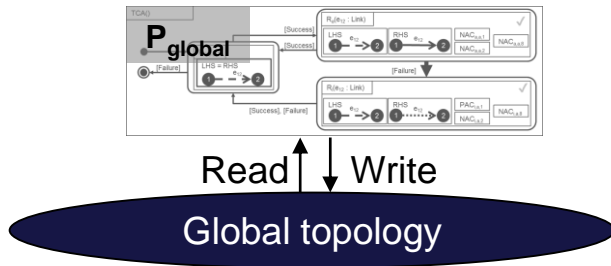
Tool support for simulation and testbed



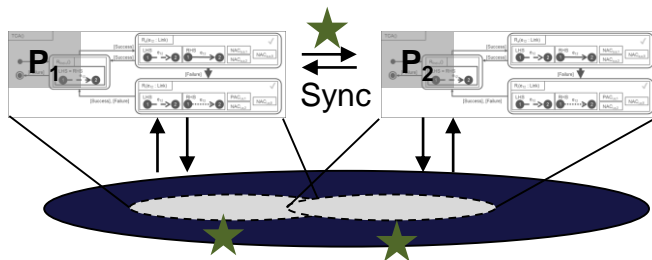
[ECMFA17] Kluge et al: cMoflon: Model-Driven Generation of Embedded C Code for Wireless Sensor Networks. In: ECMFA 2017. LNCS 10376.



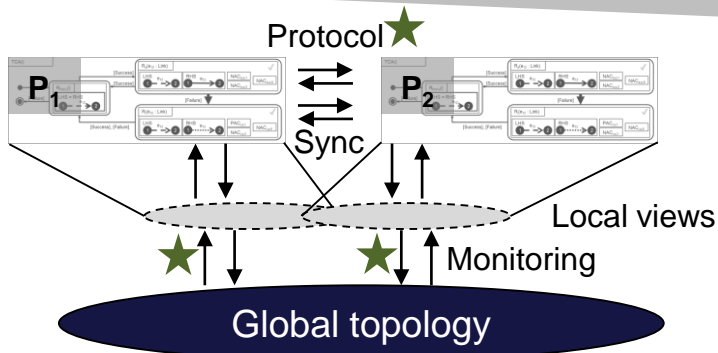
Ongoing Work: From centralized to distributed topology control algorithm specifications



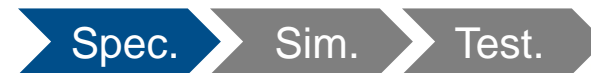
Centralized-global perspective of TC:
TC: sequential
Topology: global, consistent view



TC: distributed
Topology: local consistent view



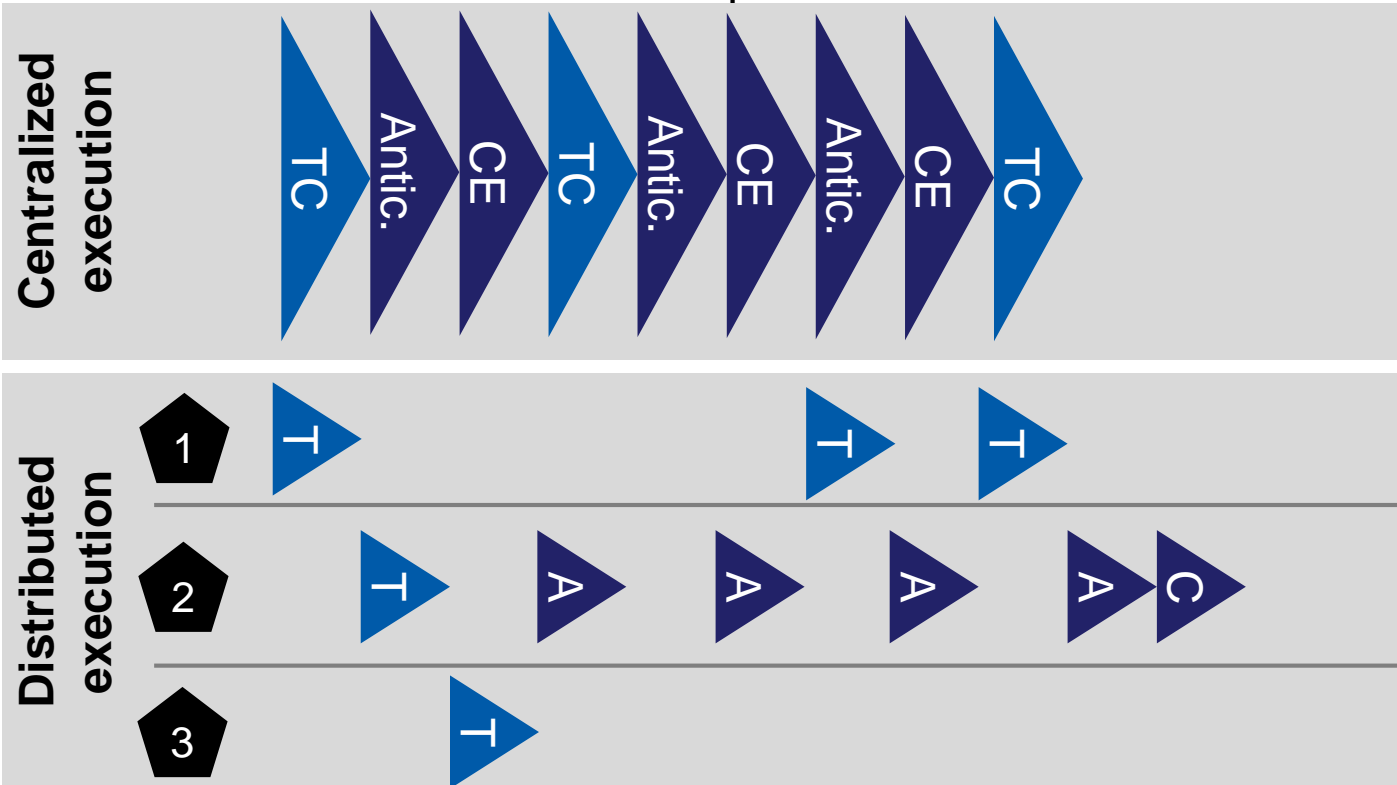
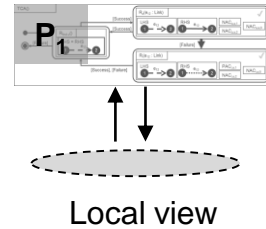
Distributed-local perspective of TC:
TC: distributed
Topology: local, **inconsistent + monitoring**



Computation model

Goal: Characterize concurrent execution + identify potential problems

- **Atomic actions:** rule applications and synchronized view of topology
- **Interleaved execution:** no two events at same point in time
- **Vertex-centric:** A node has 1 process and 1 local-view model

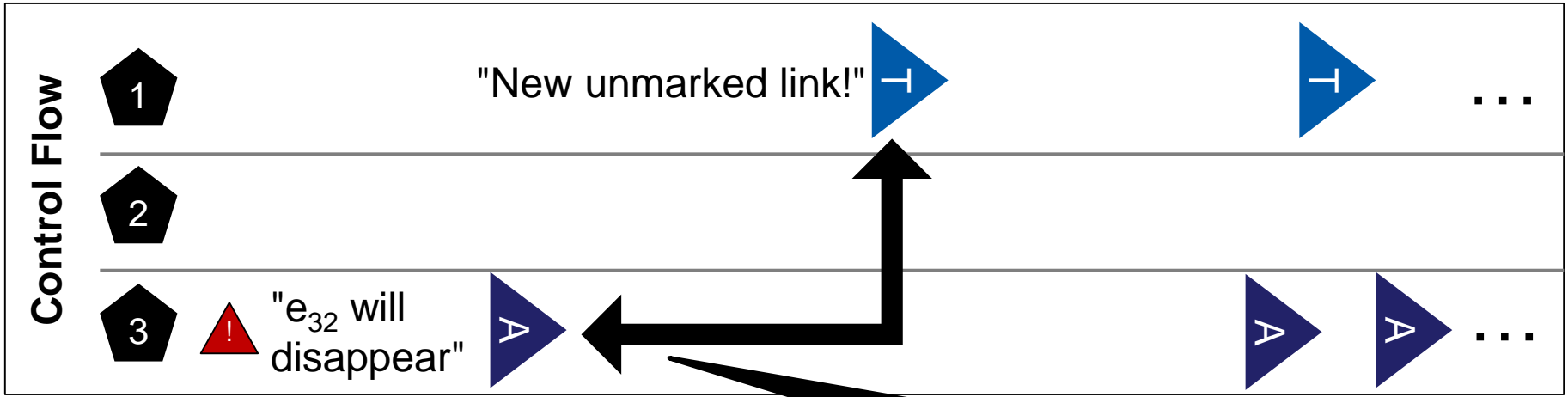
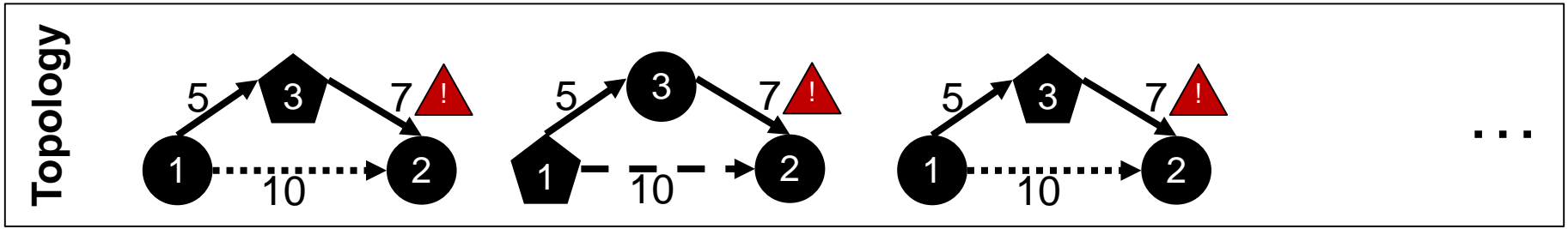


Global control flow

Per-node control flow ("happened-before")



Example: Non-termination of context handling



**Liveness problem due to Mutual dependency (necessary)
Return to same state (sufficient)**

 Pending context event

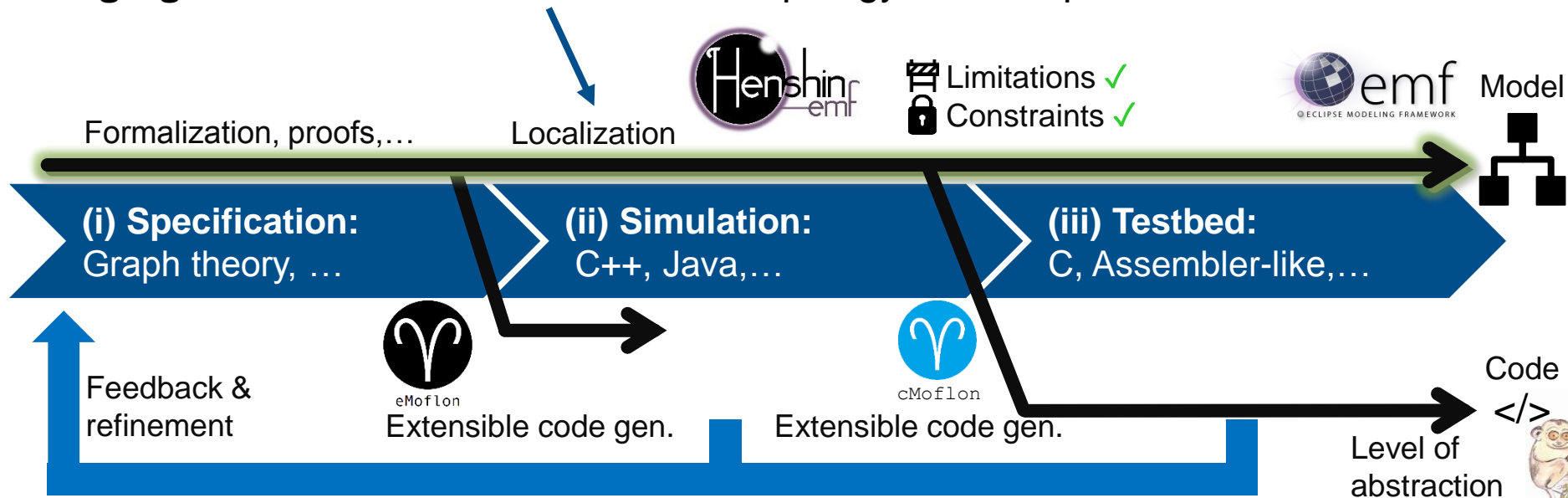


Take-Home Messages

Goal: Overcome curse of low abstraction in dev. of topology adaptations

Result: Correct-by-construction development methodology TC algorithm families

Ongoing work: Toward distributed-local topology control specifications



Supplementary material on [GitHub](#)

Thank you!



References



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Department of Electrical Engineering
and Information Technology
Real-Time Systems Lab



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